Performance Evaluation for a Hydrodynamics Application in XcalableACC PGAS Language

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Outline

• Background
• Objective
• XcalableMP and XcalableACC
• CloverLeaf in XACC
• Evaluation
• Conclusion and future work
Background (1/3)

• Clusters equipped with accelerators are increasing
  • Tianhe-2 (KNC), Piz Daint (P100), Gyoukou (PEZY-SC2), etc.
  • we refer to these clusters as “accelerated clusters”

• Programming for accelerated clusters is difficult
  • MPI for communication and CUDA / OpenCL for offloading
  • distributing data and work manually
  • communicating data using MPI function
  • managing accelerator memory using CUDA/OpenCL functions
  • describing parallel codes for accelerator
Background (2/3)

• There are some directive-based extensions to make the programming simpler

• XcalableMP (XMP) for distributed-memory systems
  • PGAS language that extends C and Fortran
  • two programming models
    • Global-view model
directive-based easy description
    • Local-view model
detailed communication using coarray

• OpenACC for accelerators
  • standard specification defined by NVIDIA, et al.
Background (3/3)

• New PGAS language XcalableACC (XACC)
  • an integration of XMP and OpenACC
  • supporting communication between accelerators
  • global-view and local-view models

• XACC has good performance and high productivity for some benchmark programs [1,2]
  • For himeno benchmark, XACC version achieves over 97% performance of MPI+OpenACC version on GPU cluster

• To assess XACC, we need evaluations in practical applications


Objective

- To evaluate performance and productivity for XACC application program
  - Target application: a hydrodynamics mini-application CloverLeaf
  - We implement CloverLeaf using XACC global-view model
  - We evaluate the performance and productivity and compare them with MPI+CUDA and MPI+OpenACC versions

http://uk-mac.github.io/CloverLeaf/
XcalableMP (XMP)

• PGAS language for distributed-memory system
  • defined by XMP Spec. W.G. of PC Cluster Consortium in Japan
  • base languages are C and Fortran
  • providing two programming models

  • Global-view model
    • directive-based programming model
      influenced by High Performance Fortran
    • providing directives for data distribution,
      work mapping, communication,
      and synchronization
    ➢ We can describe action of all nodes
      by adding directives to serial codes

  • Local-view model
    • using coarray for communication
    • not used in this work

Example of global-view model

```c
integer a(N)
!$xmp nodes p(*)
!$xmp template t(N)
!$xmp distribute t(block) onto p
!$xmp align a(i) with t(i)
!$xmp shadow a(1:1)
...

!$xmp loop (i) on t(i)
do i = 1, N
  a(i) = func(i)
end do

!$xmp reflect (a)
```

work mapping

comm.

dist.
XcalableACC (XACC)

• PGAS language for accelerated clusters
  • an integration of XMP and OpenACC

• communication between accelerators
  • Global-view model
    • acc clause for communication directives
  • Local-view model
    • coarray declaration and communication on accelerator

Example of global-view model

```plaintext
integer a(N)
!$xmp nodes p(*)
!$xmp template t(N)
!$xmp distribute t(block) onto p
!$xmp align a(i) with t(i)
!$xmp shadow a(1:1)
...
!$acc data create(a)
!$xmp loop (i) on t(i)
!$acc parallel loop
  do i = 1, N
    a(i) = func(i);
  end do
!$xmp reflect (a) acc
```

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data offload
work offload
comm.
on acc.
Example of XACC global-view model

integer a(8)
!$xmp nodes p(2)
!$xmp template t(8)
!$xmp distribute t(block) onto p
!$xmp align a(i) with t(i)

- **nodes** directive defines node set
- **template** directive defines template
  - virtual index array
  - used for data and work mapping
- **distribute** directive distributes template on nodes
- **align** directive distributes array as well as template
Example of XACC global-view model

```c
!$acc data copy (a)
!$xmp loop (i) on t(i)
!$acc kernels loop
  do i = 1, 8
    a(i) = i * 2
  end do
```

- **data directive**: allocate and copy data to accelerator memory
  - for distributed array, only assigned parts are allocated

- **xmp loop directive**: distributes loop iterations as well as the template

- **acc kernels loop directive**: offloads loop to accelerator
Example of XACC global-view model

!$xmp align a(i) with t(i)
!$xmp shadow a(1:1)
...

!$xmp reflect (a) acc

- **shadow** directive adds shadow region where can be used as halo region
- **reflect** directive updates shadow region with the value of actual region
  - acc clause specifies update shadow on accelerator memory
Omni XACC compiler

• A source–to–source XACC compiler
  • extension of Omni XMP compiler

• XACC → OpenACC + Omni XACC runtime call
  • XMP based directives are translated to runtime calls
  • enabling to use general OpenACC compilers (PGI, Cray, Omni, etc.)

• Runtime library is implemented by MPI and CUDA for GPU clusters
  • CUDA is used for pack/unpack kernels

• supporting XACC global–view model in C/Fortran
  • also XACC local–view model in C
CloverLeaf

• A hydrodynamics mini-application
  • developed by UK Atomic Weapons Establishment and University of Warwick
  • published on GitHub (https://github.com/UK-MAC/CloverLeaf)

• Solves compressible Euler equations on a 2D Cartesian grid
  • finite volume method with second-order accuracy

• Physical quantities are arranged on a staggered grid
  • often used for flow simulation on structured grid

Quantities at the corners of cells (e.g. velocity)
Quantities at the centers of cells (e.g. pressure)
CloverLeaf distributes data based on cells.

We can distribute quantities just like cells because the number of centers of cell is equal to the number of cells.
Distributed array in XACC global-view model (quantities at centers of cells)

REAL(KIND=8) pressure(x_min-2:x_max+2, y_min-2:y_max+2)
!$xmp align (i,j) with t(i,j) :: pressure
!$xmp shadow (2:2,2:2) :: pressure

distributing array according to template

adding shadow region for halo region

distribute

p(1,1)  p(2,1)

N

p(1,2)  p(2,2)
By cell based distribution, corners of cells are duplicated on divided face, but highlighted elements do not exist in original serial version.

→ We refer the duplicated parts as "extra region".

2 To allocate the extra region, we utilize shadow directive.
Distributed array in XACC global-view model
(quantities at corners of cells)

REAL(KIND=8) xvel0(x_min-2:x_max+3, y_min-2:y_max+3)
!$xmp align (i,j) with t(i,j) :: xvel0
!$xmp shadow (2:3,2:3) :: xvel0

To add the extra region, incrementing upper width of shadow region by one

• Red regions are the extra region
• We treat them as normal region in computation and don’t need to update them in halo exchange
Computation

• double nested loops construct main computation

```fortran
!$xmp loop (j,k) on t(j,k) expand(0:1,0:1)
!$acc kernels loop independent
do k=y_min,y_max+1
  !$acc loop independent
  do j=x_min,x_max+1
    xvel1(j,k) = xvel0(j,k) - ...
  enddo
enddo
```

distributing nested loops according to template same as quantity arrays

offloading nested loops and specifying iterations are data-independent

• important feature is `expand` clause

• `expand` clause expands local iteration range

To process extra region, we add `expand` clause

Normal local iteration range is same to local template range → extra region is not processed
Communication

• Main communication is halo exchange
  • Quantities at centers of cells
  ```
  xmp reflect width(2:2, 2:2) acc
  ```
  ![Diagram showing communication on accelerator]

• Quantities at corners of cells
  ```
  xmp reflect width(2:3, 2:3) acc
  ```
  ![Diagram with upper widths highlighted]

• We don’t need to update innermost shadow element, which is extra region
• However, we update shadow region including extra region because there is no clause to exclude it
Evaluation

• Three implementations
  • MPI+CUDA
  • MPI+OpenACC
  • XACC

※ We got MPI+CUDA and MPI+OpenACC versions from GitHub and use them with modifications.

• Performance evaluation
  • problem sizes are 960 × 960 and 3840 × 3840 cells
  • the number of time steps is 1000
  • execution time with strong and weak scaling

• Productivity evaluation
  • Source lines of codes (SLOC)
  • Delta SLOC (DSLOC, difference from serial version)
Evaluation environment

• HA–PACS/TCA at Center for Computational Sciences, University of Tsukuba

Node configuration and software

<p>| | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>CPU</td>
<td>Intel Xeon–E5 2680v2 2.8GHz × 2</td>
</tr>
<tr>
<td>Memory</td>
<td>DDR3 1866MHz, 128GB</td>
</tr>
<tr>
<td>GPU</td>
<td>Tesla K20X × 4</td>
</tr>
<tr>
<td>Interconnect</td>
<td>InfiniBand Mellanox Connect–X3 FDR</td>
</tr>
<tr>
<td>Software</td>
<td>PGI 16.10, CUDA 8.0, MVAPICH2 2.2, Omni Compiler 1.2.1 + extension</td>
</tr>
</tbody>
</table>

• 4 processes/node
• 1 GPU/process
• up to 64 processes on 16 nodes
Execution time (3840² cells, strong scaling)

The computational time required by the XACC version was from 4 to 12% longer than that by MPI+OpenACC.

The performance of XACC version is over 89% of MPI+CUDA and over 97% of MPI+OpenACC.

<table>
<thead>
<tr>
<th>Number of processes (X x Y)</th>
<th>Computation</th>
<th>Communication</th>
<th>MPI+OpenACC</th>
<th>XACC</th>
</tr>
</thead>
<tbody>
<tr>
<td>1x1</td>
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<tr>
<td>8x8</td>
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<td></td>
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</tr>
</tbody>
</table>

lower is better
higher is better
Computation performance degradation of XACC version

• The reason is code translation by XACC compiler
  • XACC code: Array is declared as explicit-shape array
    
    ```
    REAL(KIND=8) :: pressure(x_min-2:x_max+2, y_min-2:y_max+2)
    !$xmp align (i,j) with t(i,j) :: pressure
    ```
  
  • Translated code: Array is declared as assumed-shape or deferred-shape array and XACC runtime determines the size dynamically
    
    ```
    REAL(KIND=8) :: XMP__pressure(0:, 0:)
    ```

• array sizes are undefined at the compile time and the compiler cannot optimize the offset calculations, which increases the register utilization by GPU kernels

• it decreases the number of concurrent execution threads
Halo exchange time \((3840^2\) cells, strong scaling)\)

When the number of nodes is small, XACC version requires slightly longer time because of redundant communication.

The time becomes the same level of MPI+CUDA as the number of nodes increases and the amount of communication reduces.

When the number of nodes increases and the amount of communication reduces, the lower is better.
The computational time required by the XACC version was 4% longer than that by MPI+OpenACC.

The performance of XACC version is 89% of MPI+CUDA and 96–97% of MPI+OpenACC.
Halo exchange time \((3840^2 \text{ cells, weak scaling})\)

XACC version requires longer time at almost all configuration because the amount of redundant communication is constantly large.

Lower is better.
Proposal: extension for reflect directive

• We can specify only widths for partial shadow update
  • !$xmp reflect (array) width(2:3)

• We propose offset clause to skip inner shadow element
  • !$xmp reflect (array) width(2:2) offset(0:1)

extra region is updated

excluding the extra region

[F] !$xmp reflect ... [ offset ( reflect-offset [, reflect-offset]... ) ]
[C] #pragma xmp reflect ... [ offset ( reflect-offset [, reflect-offset]... ) ]

reflect-offset is int-expr [:int-expr]
Comparison of SLOC and DSLOC

- MPI+CUDA version requires 1.7 times more SLOC than serial version.
- XACC version requires only 1.2 times more SLOC than serial version.

- Major change in MPI+OpenACC version is non-directive code, while that of XACC version is directive addition.
  - XACC version retains a better image of the serial version than MPI+OpenACC version.

XACC version retains a better image of the serial version than MPI+OpenACC version.

Difference between MPI+OpenACC and XACC is not large.
Related work

• CloverLeaf is implemented for GPU using OpenACC, OpenCL, and CUDA [3]
  • OpenACC achieves comparable performance and high productivity comparing with OpenCL and CUDA
    • XACC also uses OpenACC for accelerator programming

• CloverLeaf is implemented using PGAS model
  OpenSHMEM and coarray [4]
  • OpenSHMEM and coarray are kind of local-view model
    • we use global-view model of XACC
  • the implementations target normal clusters
    • our work targets accelerated clusters


Conclusion

• To evaluate performance and productivity for XACC application program, we implemented a hydrodynamics mini-application CloverLeaf in XACC
  • We showed how staggered grid arrangement is implemented
• XACC has sufficient performance
  • the performance of XACC was over 89% of MPI+CUDA
• XACC has good productivity
  • XACC global-view model allows us to describe applications by adding directives to the serial version of codes

• Future work
  • implementing proposed offset clause
  • considering to improve XACC code translation